

### FlashAttention-2: Faster Attention with Better Parallelism and Work Partitioning

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- Longer sequence has proven better performance in LM and supports other modalities, code, audio.. etc.
- The main bottleneck for scaling to a longer sequence is the attention layer
- Scaling the sequence length implies a quadratic increase in runtime and memory
	- $\blacksquare$  GPT4  $\gt$  32k
	- MosiacML's MPT >> 65k
	- Anthropic Claude >> 100k

#### **Previous Work**

- FlashAttention2 is built upon the previously widely adopted FlashAttention (referred to FlashAttention1 in this context)
- Main idea is to reorder the attention computation and leverage tiling recomputation which reduces memory usage from quadratic to linear in sequence length
- It yields to 2 to 4x time speedup over the optimized baselines and up to 10 to 20x memory utilization with no approximation
- With the increased sequence length FlashAttention underperform other primitives such as GEMM (General Matrix-multiply)
- Forward pass reaches 30-50% of theoretical maximum FLOPs/s
- Backward pass is even more challenging, reaching only 25-35% of the maximum throughput
- Optimized GEMM can reach up to 80-90% of the theoretical maximum device throughput



#### **Background – Hardware Characteristics**



- **Memory Hierarchy** comprises of high bandwidth memory (HBM) [Slow and Large], and onchip SRAM (aka shared memory) [Fast and Small]
- **Execution Model** Kernal is the number of threads
- Threads are organized into blocks scheduled to run on streaming multiprocessors (SMs)
- In each thread block, threads are grouped into warps (32 threads)
- Threads within a warp threads can communicate faster
- Warps within a thread block can communicate by reading from / writing to shared memory
- Each kernel loads inputs from HBM to registers and SRAM, computes, then writes outputs to HBM.

#### **Background – Standard Attention Implementation**



 $\mathbf{S} = \mathbf{Q}\mathbf{K}^{\top} \in \mathbb{R}^{N \times N}, \quad \mathbf{P} = \text{softmax}(\mathbf{S}) \in \mathbb{R}^{N \times N}, \quad \mathbf{O} = \mathbf{P}\mathbf{V} \in \mathbb{R}^{N \times d}$ 

The Softmax is applied row-wise.

#### **Standard attention implementation**

- 1. Calls GEMM (Matrix-Multiply) to compute **S**
- 2. Writes results to HBM
- 3. Load from HBM to compute the Softmax
- 4. Rewrite **P** to HBM
- 5. Calls GEMM to compute **O**

This frequent memory access cause s slow execution and requires O(N<sup>2</sup> ) since **S** and **P** are materialize P has to be saved for backward pass to compute the gradients

#### **Background – Flash Attention or FA1**



- Forward Pass: Utilize a classical technique of tiling to reduce memory IOs, by:
- 1. loading blocks of inputs from HBM to SRAM
- 2. computing attention for that block, then
- 3. updating the output without writing the large intermediate matrices **S** and **P** to HBM
- Using online Softmax to split the attention computation into blocks, rescale the output of each block to get non-approximate results.
- Reduces the number of memory's reads/writes
- FlashAttention yields 2-4× wall-clock speedup over-optimized baseline attention implementations.

## **Background – Standard Softmax**



#### ■ Considering two blocks in S

$$
m = \max(\text{rowmax}(\mathbf{S}^{(1)}), \text{rowmax}(\mathbf{S}^{(2)})) \in \mathbb{R}^{B_r}
$$
  
\n
$$
\ell = \text{rowsum}(e^{\mathbf{S}^{(1)} - m}) + \text{rowsum}(e^{\mathbf{S}^{(2)} - m}) \in \mathbb{R}^{B_r}
$$
  
\n
$$
\mathbf{P} = [\mathbf{P}^{(1)} \quad \mathbf{P}^{(2)}] = \text{diag}(\ell)^{-1} \left[e^{\mathbf{S}^{(1)} - m} e^{\mathbf{S}^{(2)} - m}\right] \in \mathbb{R}^{B_r \times 2B_c}
$$
  
\n
$$
\mathbf{O} = [\mathbf{P}^{(1)} \quad \mathbf{P}^{(2)}] \left[\mathbf{V}^{(1)}\right] = \text{diag}(\ell)^{-1} e^{\mathbf{S}^{(1)} - m} \mathbf{V}^{(1)} + e^{\mathbf{S}^{(2)} - m} \mathbf{V}^{(2)} \in \mathbb{R}^{B_r \times d}
$$

-m for stabilization

# **Background – FA1**



Online softmax instead computes "local" softmax with respect to each block and rescale to get the right output at the end

$$
m^{(1)} = \text{rowmax}(\mathbf{S}^{(1)}) \in \mathbb{R}^{B_r}
$$
  
\n
$$
\ell^{(1)} = \text{rowsum}(e^{\mathbf{S}^{(1)} - m^{(1)}}) \in \mathbb{R}^{B_r}
$$
  
\n
$$
\tilde{\mathbf{P}}^{(1)} = \text{diag}(\ell^{(1)})^{-1}e^{\mathbf{S}^{(1)} - m^{(1)}} \in \mathbb{R}^{B_r \times B_c}
$$
  
\n
$$
\mathbf{O}^{(1)} = \tilde{\mathbf{P}}^{(1)}\mathbf{V}^{(1)} = \text{diag}(\ell^{(1)})^{-1}e^{\mathbf{S}^{(1)} - m^{(1)}}\mathbf{V}^{(1)} \in \mathbb{R}^{B_r \times d}
$$
  
\n
$$
m^{(2)} = \max(m^{(1)}, \text{rowmax}(\mathbf{S}^{(2)})) = m
$$
  
\n
$$
\ell^{(2)} = e^{m^{(1)} - m^{(2)}}\ell^{(1)} + \text{rowsum}(e^{\mathbf{S}^{(2)} - m^{(2)}}) = \text{rowsum}(e^{\mathbf{S}^{(1)} - m}) + \text{rowsum}(e^{\mathbf{S}^{(2)} - m}) = \ell
$$
  
\n
$$
\tilde{\mathbf{P}}^{(2)} = \text{diag}(\ell^{(2)})^{-1}e^{\mathbf{S}^{(2)} - m^{(2)}}
$$
  
\n
$$
\mathbf{O}^{(2)} = \text{diag}(\ell^{(1)}/\ell^{(2)})^{-1}\mathbf{O}^{(1)} + \tilde{\mathbf{P}}^{(2)}\mathbf{V}^{(2)} = \text{diag}(\ell^{(2)})^{-1}e^{\mathbf{S}^{(1)} - m}\mathbf{V}^{(1)} + \text{diag}(\ell^{(2)})^{-1}e^{\mathbf{S}^{(2)} - m}\mathbf{V}^{(2)} = \mathbf{O}.
$$



### **(1) Reduce Non-Matmul FLOPs Operations in FA2**

$$
m^{(1)} = \text{rowmax}(\mathbf{S}^{(1)}) \in \mathbb{R}^{B_r}
$$
  
\n
$$
\ell^{(1)} = \text{rowsum}(e^{\mathbf{S}^{(1)} - m^{(1)}}) \in \mathbb{R}^{B_r}
$$
  
\n
$$
\mathbf{O}^{(1)} = e^{\mathbf{S}^{(1)} - m^{(1)}} \mathbf{V}^{(1)} \in \mathbb{R}^{B_r \times d}
$$
  
\n
$$
m^{(2)} = \max(m^{(1)}, \text{rowmax}(\mathbf{S}^{(2)})) = m
$$
  
\n
$$
\ell^{(2)} = e^{m^{(1)} - m^{(2)}} \ell^{(1)} + \text{rowsum}(e^{\mathbf{S}^{(2)} - m^{(2)}}) = \text{rowsum}(e^{\mathbf{S}^{(1)} - m}) + \text{rowsum}(e^{\mathbf{S}^{(2)} - m}) = \ell
$$
  
\n
$$
\tilde{\mathbf{P}}^{(2)} = \text{diag}(\ell^{(2)})^{-1} e^{\mathbf{S}^{(2)} - m^{(2)}}
$$
  
\n
$$
\tilde{\mathbf{O}}^{(2)} = \text{diag}(e^{m^{(1)} - m^{(2)}}) \tilde{\mathbf{O}}^{(1)} + e^{\mathbf{S}^{(2)} - m^{(2)}} \mathbf{V}^{(2)} = e^{s^{(1)} - m} \mathbf{V}^{(1)} + e^{s^{(2)} - m} \mathbf{V}^{(2)}
$$
  
\n
$$
\mathbf{O}^{(2)} = \text{diag}(\ell^{(2)})^{-1} \tilde{\mathbf{O}}^{(2)} = \mathbf{O}.
$$

# **(2) Parallelizing Attention Computation**

- $FA1$  parallelizes over batch size (B) and the number of heads >> running many threads blocks together
- Each thread runs on a single streaming multiprocessor (SM) e.g., A100 has 108 SMs
- Unfortunately, with long sequence s, we cannot afford large batch size s or large number s of heads
- This leads to many idle SMs
- FA2 parallelizes additionally over sequence length, so supports three parallelization





#### **Backward pass**





### **(2) Parallelizing Attention Computation**



- FA1 works in two loops
- **EX** First over the j-th(k, v) blocks, and the second runs over the i-th Q
- **They compute in SRAM (fast and small), and update O<sup>i</sup>,**  $I_i$ and  $m_i$  in HBM (large and slow)
- In FA2 It swaps the order of the loop
- Placing  $Q_i$  outside  $\geq$  the loop goes through different blocks of the Q matrix
- Inner operates on K and V
- Swapping the order offers sequence length parallelization as the #  $\overline{Q}$  == to the length of the sequence
- In autoregressive attention >> set the upper triangle to infinity and leave them uncomputed

Forward pass



#### Backward pass



### **(3) Work Partitioning between Warps**



- Typically use 4 or 8 warps per thread block
- Every warps contains 32 threads
- In FA1 "Split-K" technique:
	- K and V are split across 4 warps, while Q is kept accessible by all warps
	- Each warp multiplies to get a slice of  $QK^{T}$
	- $\blacksquare$  Then multiply with a slice of V



#### **(3) Work Partitioning between Warps**



#### In  $FA2$ :

- we split  $Q$  across 4 warps
- While keeping K and V accessible by all warps
- After each warp computes a slice of QK<sup>T</sup> , they only need to multiply with their shared slice of V, to get their corresponding slice of the output O
- No need to communicate between warps
- Reduction in shared memory's reads/writes >> speedup
- The same applies to backward pass





203

201

9795

 $8k$ 

182

 $8k$ 

 $16k$ 

189

92

 $16k$ 

196

 $4k$ 

173

87

 $4k$ 

Sequence length

155

187

 $2k$ 

 $1k$ 

133

1k

#### **Results**

- A100
- Different Head Dimension [64, 128]
- W/O Causal mask



(c) With causal mask, head dimension 64



Sequence length

 $2k$ 

Figure 4: Attention forward  $+$  backward speed on A100 GPU

### **Results**

- Forward Speed on A100
- Different Head Dimension [64, 128]
- W/O Causal mask



(c) With causal mask, head dimension 64

(d) With causal mask, head dimension 128

 $2k$ 

Sequence length

 $4k$ 

 $8k$ 

 $1k$ 

Attention forward speed (A100 80GB SXM4)

 $1k$ 

132

 $2k$ 

Sequence length

Attention forward speed (A100 80GB SXM4)

187

 $4k$ 

198

8k

200

222

224

223

163

 $16k$ 

197

 $16k$ 

Figure 5: Attention forward speed on A100 GPU



### **Results**

- Backward Speed on A100
- Different Head Dimension [64, 128]
- W/O Causal mask



Attention backward speed (A100 80GB SXM4)

 $\blacksquare$  Pytorch FlashAttention

xformers

200

(c) With causal mask, head dimension 64

Sequence length



(b) Without causal mask, head dimension 128





Figure 6: Attention backward speed on A100 GPU

#### **Results – Training Speed**



Table 1: Training speed (TFLOPs/s/GPU) of GPT-style models on 8×A100 GPUs. FLASHATTENTION-2 reaches up to 225 TFLOPs/s (72% model FLOPs utilization). We compare against a baseline running without FLASHATTENTION.



#### **H100 GPU**

- Different Head Dimension [64, 128]
- W/O Causal mask





Sequence length



(d) With causal mask, head dimension 128

Figure 7: Attention forward + backward speed on H100 GPU



(b) Without causal mask, head dimension 128

Attention forward + backward speed (H100 80GB SXM5)



# Thank you